[MS-XCA]: Xpress Compression Algorithm

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Errata below are for Protocol Document Version V5.0 – 2018/09/12.

Errata Published*	Description

2020/02/17	In Section 2.3.4, Processing, we updated the pseudocode for the encoding method for match lengths greater than 65535.
	
	Changed from:
	
	`If MatchLength >= 7`
	`  MatchLength -= 7`
	`  If LastLengthHalfByte == 0`
	`  LastLengthHalfByte = OutputPosition`
	`  Write the byte value min(MatchLength, 15) to OutputPosition`
	`  OutputPosition += 1`
	`  Else`
	`  OutputBuffer[LastLengthHalfByte] |= min(15, MatchLength) << 4`
	`  LastLengthHalfByte = 0`
	`  If MatchLength >= 15`
	`  MatchLength -= 15`
	`  Write the byte value min(MatchLength, 255) to OutputPosition`
	`  OutputPosition += 1`
	`  If MatchLength >= 255`
	`  MatchLength += 15 + 7`
	`  Write the 2-byte value MatchLength to OutputPosition`
	`  OutputPosition += 2`
	
	Changed to:
	
	`If MatchLength < 7`
	`  // This is the simple case. The length fits in 3 bits.`

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	MatchOffset += MatchLength
	Write MatchOffset the 2-byte value to OutputPosition
	OutputPosition += 2
	Else
	// The length does not fit 3 bits. Record a special value to
	// indicate a longer length.
	MatchOffset |= 7
	Write MatchOffset the 2-byte value to OutputPosition
	OutputPosition += 2
	MatchLength -= 7
	// Try to encode the length in the next 4 bits. If we previously
	// encoded a 4-bit length, we'll use the high 4 bits from that byte.
	If LastLengthHalfByte == 0
	LastLengthHalfByte = OutputPosition
	If MatchLength < 15
	Write single byte value of MatchLength to OutputPosition
	OutputPosition += 1
	Else
	Write single byte value of 15 to OutputPosition
	OutputPosition++
	goto EncodeExtraLen
	Else
	If MatchLength < 15
	OutputBuffer[LastLengthHalfByte] |= MatchLength << 4
	LastLengthHalfByte = 0
	Else
	OutputBuffer[LastLengthHalfByte] |= 15 << 4
	LastLengthHalfByte = 0
	EncodeExtraLen:
	// We've already used 3 bits + 4 bits to encode the length
	// Next use the next byte.
	MatchLength -= 15
	If MatchLength < 255
	Write single byte value of MatchLength to OutputPosition

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	OutputPosition += 1
	Else
	// Use two more bytes for the length
	Write single byte value of 255 to OutputPosition
	OutputPosition += 1
	MatchLength += 7 + 15
	If MatchLength < (1 << 16)
	Write two-byte value MatchLength to OutputPosition
	OutputPosition += 2
	Else
	Write two-byte value of 0 to OutputPosition
	OutputPosition += 2
	Write four-byte value of MatchLength to OutputPosition
	OutputPosition += 4

2020/02/17	In Section 2.3.4 Processing, we added clarifying information about the maximum MatchLength. Changed from: The fastest variant of the Xpress Compression Algorithm avoids the cost of the Huffman[IEEE-MRC] pass by encoding the LZ77 [UASDC] literals and matches in a simple way. The encoding process is similar to the method described in section 2.1.4.1, with the key difference that the largest match offset it can encode is 8192 instead of the 65535 limit of the Huffman format. The literal or match flags are encoded in 32-bit chunks. Literals are encoded with a simple byte value. Matches are encoded with a 16-bit value, where the high 13 bits represent the offset and the low 3 bits represent the length. Long lengths are encoded with an additional 4 bits, then 8 bits, and then 16 bits. The following pseudocode provides an outline of the encoding method. Changed to: The fastest variant of the Xpress Compression Algorithm avoids the cost of the Huffman[IEEE-MRC] pass by encoding the LZ77 [UASDC] literals and matches in a simple way. The encoding process is similar to the method described in section 2.1.4.1, with the key difference that the largest match offset it can encode is 8192 instead of the 65535 limit of the Huffman format. The literal or match flags are encoded in 32-bit chunks. Literals are encoded with a simple byte value. Matches are encoded with a 16-bit value, where the high 13 bits represent the offset and the low 3 bits represent the length. Long lengths are encoded with an additional 4 bits, then 8 bits, and then 16 bits. The MatchLength is represented by a ULONG, a 32-bit unsigned integer (see [MS-DTYP] section 2.2.51); therefore, the maximum value is 4,294,967,295. The following pseudocode provides an outline of the encoding method.
2020/02/17	In Section 2.2.4 Processing, we corrected the pseudocode by replacing DecodedValue with HuffmanSymbol and added a clarifying comment to the pseudocode to explain why the HuffmanSymbol needs to be right-shifted by 4 bits.

Errata Published*	Description

	Changed from:
	
	…
	Loop until a decompression terminating condition
	Build the decoding table
	CurrentPosition = 256 // start at the end of the Huffman table
	NextBits = Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	NextBits <<= 16
	NextBits |= Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	ExtraBits = 16
	BlockEnd = OutputPosition + 65536
	Loop until a block terminating condition
	If OutputPosition >= BlockEnd then terminate block processing
	Loop until a literal processing terminating condition
	Next15Bits = NextBits >> (32 – 15)
	HuffmanSymbol = DecodingTable[Next15Bits]
	HuffmanSymbolBitLength = the bit length of HuffmanSymbol, from the table in
	the input buffer
	If HuffmanSymbol <= 0
	NextBits <<= HuffmanSymbolBitLength
	ExtraBits -= HuffmanSymbolBitLength
	Do
	HuffmanSymbol = - HuffmanSymbol
	HuffmanSymbol += (NextBits >> 31)
	NextBits *= 2
	ExtraBits = ExtraBits - 1
	HuffmanSymbol = DecodingTable[HuffmanSymbol]
	While DecodedValue <= 0
	Else
	DecodedBitCount = DecodedValue & 15
	NextBits <<= DecodedBitCount
	ExtraBits -= DedcodedBitCount

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	HuffmanSymbol >>= 4 HuffmanSymbol -= 256 If ExtraBits < 0 NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits) ExtraBits += 16 CurrentPosition += 2 If HuffmanSymbol >= 0 If HuffmanSymbol == 0 If the entire input buffer has been read and the expected decompressed size has been written to the output buffer Decompression is complete. Return with success. Terminate literal processing Else Output the byte value of HuffmanSymbol to the output stream End of literal processing Loop MatchLength = HuffmanSymbol mod 16 MatchOffsetBitLength = HuffmanSymbol / 16 If MatchLength == 15 MatchLength = ReadByte(InputBuffer + CurrentPosition) CurrentPosition += 1 If MatchLength == 255 MatchLength = Read16Bits(InputBuffer + CurrentPosition) CurrentPosition += 2 If MatchLength < 15 The compressed data is invalid. Return error. MatchLength = MatchLength - 15 MatchLength = MatchLength + 15 MatchLength = MatchLength + 3 MatchOffset = NextBits >> (32 – MatchOffsetBitLength) MatchOffset += (1 << MatchOffsetBitLength) NextBits <<= MatchOffsetBitLength ExtraBits -= MatchOffsetBitLength If ExtraBits < 0

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	NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits)
	ExtraBits += 16
	CurrentPosition += 2
	For i = 0 to MatchLength - 1
	Output OutputBuffer[CurrentOutputPosition – MatchOffset + i]
	End of block loop
	End of decoding loop
	
	Changed to:
	
	…
	Loop until a decompression terminating condition
	Build the decoding table
	CurrentPosition = 256 // start at the end of the Huffman table
	NextBits = Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	NextBits <<= 16
	NextBits |= Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	ExtraBits = 16
	BlockEnd = OutputPosition + 65536
	Loop until a block terminating condition
	If OutputPosition >= BlockEnd then terminate block processing
	Loop until a literal processing terminating condition
	Next15Bits = NextBits >> (32 – 15)
	HuffmanSymbol = DecodingTable[Next15Bits]
	HuffmanSymbolBitLength = the bit length of HuffmanSymbol, from the table in
	the input buffer
	If HuffmanSymbol <= 0
	NextBits <<= HuffmanSymbolBitLength
	ExtraBits -= HuffmanSymbolBitLength
	Do
	HuffmanSymbol = - HuffmanSymbol

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	HuffmanSymbol += (NextBits >> 31) NextBits *= 2 ExtraBits = ExtraBits - 1 HuffmanSymbol = DecodingTable[HuffmanSymbol] **While HuffmanSymbol <= 0** Else **DecodedBitCount = HuffmanSymbol & 15** NextBits <<= DecodedBitCount ExtraBits -= DedcodedBitCount HuffmanSymbol >>= 4 // Shift by 4 bits to get the symbol value // (the lower 4 bits are the bit length of the symbol) HuffmanSymbol -= 256 If ExtraBits < 0 NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits) ExtraBits += 16 CurrentPosition += 2 If HuffmanSymbol >= 0 If HuffmanSymbol == 0 If the entire input buffer has been read and the expected decompressed size has been written to the output buffer Decompression is complete. Return with success. Terminate literal processing Else Output the byte value of HuffmanSymbol to the output stream End of literal processing Loop MatchLength = HuffmanSymbol mod 16 MatchOffsetBitLength = HuffmanSymbol / 16 If MatchLength == 15 MatchLength = ReadByte(InputBuffer + CurrentPosition) CurrentPosition += 1 If MatchLength == 255 MatchLength = Read16Bits(InputBuffer + CurrentPosition)

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	CurrentPosition += 2 If MatchLength < 15 The compressed data is invalid. Return error. MatchLength = MatchLength - 15 MatchLength = MatchLength + 15 MatchLength = MatchLength + 3 MatchOffset = NextBits >> (32 – MatchOffsetBitLength) MatchOffset += (1 << MatchOffsetBitLength) NextBits <<= MatchOffsetBitLength ExtraBits -= MatchOffsetBitLength If ExtraBits < 0 NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits) ExtraBits += 16 CurrentPosition += 2 For i = 0 to MatchLength - 1 Output OutputBuffer[CurrentOutputPosition – MatchOffset + i] End of block loop End of decoding loop
2019/12/09	In Section 2.1, LZ77+Huffman Compression Algorithm Details, described how data is processed for the Huffman variant. Changed from: The overall compression algorithm for the Huffman [IEEE-MRC] variant can be divided into three stages, which are performed in this order: ... Changed to: The overall compression algorithm for the Huffman [IEEE-MRC] variant can handle an arbitrary amount of data. Data is processed in 64k blocks, and the encoded results are stored in-order. After the final block, the end-of-file (EOF) symbol is encoded. Each 64k block is run through three stages, which are performed in this order: ... In Section 2.2.4, Processing, described the decompression process and clarified how the compression stream handles the bytes for long match lengths in the pseudocode. Changed from:

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	The decompression algorithm uses the 256-byte Huffman table to reconstruct the canonical Huffman [IEEE-MRC] representations of each symbol. Next, the Huffman stream of LZ77 ([UASDC]) literals and matches is decoded to reproduce the original data.
	
	The following method can be used to construct a decoding table. The decoding table will have 2^15 entries because 15 is the maximum bit length permitted by the Xpress Compression Algorithm for a Huffman code. If a symbol has a bit length of X, it has 2^(15 – X) entries in the table that point to its value. The order of symbols in the table is sorted by bit length (from low to high), and then by symbol value (from low to high). These requirements represent the agreement of canonicalness with the compression end of the algorithm. The following pseudocode shows the table construction method:
	...
	The compression stream is designed to be read in (mostly) 16-bit chunks, with a 32-bit register maintaining at least the next 16 bits of input. This strategy allows the code to seamlessly handle the bytes for long match lengths, which would otherwise be awkward. The following pseudocode demonstrates this method.
	
	Build the decoding table
	CurrentPosition = 256 // start at the end of the Huffman table
	NextBits = Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	NextBits <<= 16
	NextBits |= Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	ExtraBits = 16
	Loop until a terminating condition
	Next15Bits = NextBits >> (32 – 15)
	HuffmanSymbol = DecodingTable[Next15Bits]
	HuffmanSymbolBitLength = the bit length of HuffmanSymbol, from the table in
	the input buffer
	NextBits <<= HuffmanSymbolBitLength
	ExtraBits -= HuffmanSymbolBitLength
	If ExtraBits < 0
	NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits)
	ExtraBits += 16
	CurrentPosition += 2
	If HuffmanSymbol < 256
	Output the byte value HuffmanSymbol to the output stream.
	Else If HuffmanSymbol == 256 and
	the entire input buffer has been read and
	the expected decompressed size has been written to the output buffer
	Decompression is complete. Return with success.
	Else
	HuffmanSymbol = HuffmanSymbol - 256
	MatchLength = HuffmanSymbol mod 16
	MatchOffsetBitLength = HuffmanSymbol / 16
	If MatchLength == 15
	MatchLength = ReadByte(InputBuffer + CurrentPosition)
	CurrentPosition += 1
	If MatchLength == 255

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	MatchLength = Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	If MatchLength < 15
	The compressed data is invalid. Return error.
	MatchLength = MatchLength - 15
	MatchLength = MatchLength + 15
	MatchLength = MatchLength + 3
	MatchOffset = NextBits >> (32 – MatchOffsetBitLength)
	MatchOffset += (1 << MatchOffsetBitLength)
	NextBits <<= MatchOffsetBitLength
	ExtraBits -= MatchOffsetBitLength
	If ExtraBits < 0
	Read the next 2 bytes the same as the preceding (ExtraBits < 0) case
	For i = 0 to MatchLength - 1
	Output OutputBuffer[CurrentOutputPosition – MatchOffset + i]
	...
	
	Changed to:
	The decompression processes a series of blocks to form the decompressed output. Each block is processed in-order, and its decoded content written to the output stream is in-order. When processing a block, we check for terminating conditions for both block and overall decoding.
	
	The decompression algorithm uses the 256-byte Huffman table to reconstruct the canonical Huffman [IEEE-MRC] representations of each symbol. Next, the Huffman stream of LZ77 ([UASDC]) literals and matches is decoded to reproduce the original data.
	
	The following method can be used to construct a decoding table. The decoding table will have 2^15 entries because 15 is the maximum bit length permitted by the Xpress Compression Algorithm for a Huffman code. If a symbol has a bit length of X, it has 2^(15 – X) entries in the table that point to its value. The order of symbols in the table is sorted by bit length (from low to high), and then by symbol value (from low to high). These requirements represent the agreement of canonicalness with the compression end of the algorithm. The following pseudocode shows the table construction method:
	...
	The compression stream is designed to be read in (mostly) 16-bit chunks, with a 32-bit register maintaining at least the next 16 bits of input. This strategy allows the code to seamlessly handle the bytes for long match lengths, which would otherwise be awkward. The following pseudocode demonstrates this method.
	
	Loop until a decompression terminating condition
	Build the decoding table
	CurrentPosition = 256 // start at the end of the Huffman table
	NextBits = Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	NextBits <<= 16
	NextBits |= Read16Bits(InputBuffer + CurrentPosition)
	CurrentPosition += 2
	ExtraBits = 16
	BlockEnd = OutputPosition + 65536
	
	Loop until a block terminating condition

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	If OutputPosition >= BlockEnd then terminate block processing Loop until a literal processing terminating condition Next15Bits = NextBits >> (32 – 15) HuffmanSymbol = DecodingTable[Next15Bits] HuffmanSymbolBitLength = the bit length of HuffmanSymbol, from the table in the input buffer If HuffmanSymbol <= 0 NextBits <<= HuffmanSymbolBitLength ExtraBits -= HuffmanSymbolBitLength Do HuffmanSymbol = - HuffmanSymbol HuffmanSymbol += (NextBits >> 31) NextBits *= 2 ExtraBits = ExtraBits - 1 HuffmanSymbol = DecodingTable[HuffmanSymbol] While DecodedValue <= 0 Else DecodedBitCount = DecodedValue & 15 NextBits <<= DecodedBitCount ExtraBits -= DedcodedBitCount HuffmanSymbol >>= 4 HuffmanSymbol -= 256 If ExtraBits < 0 NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits) ExtraBits += 16 CurrentPosition += 2 If HuffmanSymbol >= 0 If HuffmanSymbol == 0 If the entire input buffer has been read and the expected decompressed size has been written to the output buffer Decompression is complete. Return with success. Terminate literal processing Else Output the byte value of HuffmanSymbol to the output stream End of literal processing Loop MatchLength = HuffmanSymbol mod 16 MatchOffsetBitLength = HuffmanSymbol / 16 If MatchLength == 15 MatchLength = ReadByte(InputBuffer + CurrentPosition) CurrentPosition += 1 If MatchLength == 255 MatchLength = Read16Bits(InputBuffer + CurrentPosition) CurrentPosition += 2 If MatchLength < 15 The compressed data is invalid. Return error. MatchLength = MatchLength - 15

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	MatchLength = MatchLength + 15 MatchLength = MatchLength + 3 MatchOffset = NextBits >> (32 – MatchOffsetBitLength) MatchOffset += (1 << MatchOffsetBitLength) NextBits <<= MatchOffsetBitLength ExtraBits -= MatchOffsetBitLength If ExtraBits < 0 NextBits |= Read16Bits(InputBuffer + CurrentPosition) << (-ExtraBits) ExtraBits += 16 CurrentPosition += 2 For i = 0 to MatchLength - 1 Output OutputBuffer[CurrentOutputPosition – MatchOffset + i] End of block loop End of decoding loop ...

2019/09/02	In Section 2.4.4, Processing, pseudocode supporting longer matches has been updated
	
	Changed from:
	...
	The match length can be greater than the match offset, and this necessitates the 1-byte-at-a-time copying strategy shown in the following pseudocode.
	
	`BufferedFlags = 0`
	`BufferedFlagCount = 0`
	`InputPosition = 0`
	`OutputPosition = 0`
	`LastLengthHalfByte = 0`
	`Loop until break instruction or error`
	`If BufferedFlagCount == 0`
	`BufferedFlags = read 4 bytes at InputPosition`
	`InputPosition += 4`
	`BufferedFlagCount = 32`
	`BufferedFlagCount = BufferedFlagCount – 1`
	`If (BufferedFlags & (1 << BufferedFlagCount)) == 0`
	`Copy 1 byte from InputPosition to OutputPosition.  Advance both.`
	`Else`
	`If InputPosition == InputBufferSize`
	`Decompression is complete.  Return with success.`
	`MatchBytes = read 2 bytes from InputPosition`
	`InputPosition += 2`
	`MatchLength = MatchBytes mod 8`
	`MatchOffset = (MatchBytes / 8) + 1`
	`If MatchLength == 7`
	`If LastLengthHalfByte == 0`
	`MatchLength = read 1 byte from InputPosition`
	`MatchLength = MatchLength mod 16`
	`LastLengthHalfByte = InputPosition`
	`InputPosition += 1`
	`Else`
	`MatchLength = read 1 byte from LastLengthHalfByte position`
	`MatchLength = MatchLength / 16`
	`LastLengthHalfByte = 0`
	`If MatchLength == 15`
	`MatchLength = read 1 byte from InputPosition`
	`InputPosition += 1`
	`If MatchLength == 255`

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```
                        MatchLength = read 2 bytes from InputPosition
                        InputPosition += 2
                        If MatchLength < 15 + 7
                           Return error.
                        MatchLength -= (15 + 7)
                    MatchLength += 15
                MatchLength += 7
            MatchLength += 3
            For i = 0 to MatchLength – 1
                Copy 1 byte from OutputBuffer[OutputPosition – MatchOffset]
                OutputPosition += 1
```

Changed to:

...

The match length can be greater than the match offset, and this necessitates the 1-byte-at-a-time copying strategy shown in the following pseudocode.

```
BufferedFlags = 0
BufferedFlagCount = 0
InputPosition = 0
OutputPosition = 0
LastLengthHalfByte = 0
Loop until break instruction or error
    If BufferedFlagCount == 0
        BufferedFlags = read 4 bytes at InputPosition
        InputPosition += 4
        BufferedFlagCount = 32
    BufferedFlagCount = BufferedFlagCount – 1
    If (BufferedFlags & (1 << BufferedFlagCount)) == 0
        Copy 1 byte from InputPosition to OutputPosition.  Advance both.
    Else
        If InputPosition == InputBufferSize
            Decompression is complete.  Return with success.
        MatchBytes = read 2 bytes from InputPosition
        InputPosition += 2
        MatchLength = MatchBytes mod 8
        MatchOffset = (MatchBytes / 8) + 1
        If MatchLength == 7
            If LastLengthHalfByte == 0
                MatchLength = read 1 byte from InputPosition
                MatchLength = MatchLength mod 16
                LastLengthHalfByte = InputPosition
                InputPosition += 1
            Else
                MatchLength = read 1 byte from LastLengthHalfByte position
                MatchLength = MatchLength / 16
                LastLengthHalfByte = 0
            If MatchLength == 15
                MatchLength = read 1 byte from InputPosition
                InputPosition += 1
                If MatchLength == 255
                    MatchLength = read 2 bytes from InputPosition
                    InputPosition += 2
                    If MatchLength == 0
                        MatchLength = read 4 bytes from InputPosition
                        InputPosition += 4 bytes
                    If MatchLength < 15 + 7
                       Return error.
                    MatchLength -= (15 + 7)
                MatchLength += 15
            MatchLength += 7
        MatchLength += 3
```

Errata Published*	Description
	For i = 0 to MatchLength – 1 Copy 1 byte from OutputBuffer[OutputPosition – MatchOffset] OutputPosition += 1
2019/07/08	In Section 2.1.4.2, Huffman Code Construction Phase, clarified that the sorting algorithm used in the Huffman Code construction phase is stable. Changed from: ... The following flowchart illustrates the length-limited canonical Huffman code construction method. ... Changed to: ... The following flowchart illustrates the length-limited canonical Huffman code construction method. Note that the sorting algorithm used in the Huffman Code construction phase is stable. ...
2019/07/08	In Section 2.1.4.3 Final Encoding Phase, clarified that some implementations of the decompression algorithm expect a terminating Huffman symbol and that it is recommended the encoding algorithm append this symbol. Changed from: Some implementations of the decompression algorithm expect an extra symbol to mark the end of the data. For example, certain implementations fail during decompression if the Huffman symbol 256 is not found after the actual data. For this reason, the encoding algorithm appends this symbol and increments the count of symbol 256 before the Huffman codes are constructed. Changed to: Implementations of the decompression algorithm may expect an extra symbol to mark the end of the data. For example, certain implementations fail during decompression if the Huffman symbol 256 is not found after the actual data. For this reason, the encoding algorithm SHOULD append this symbol and increment the count of symbol 256 before the Huffman codes are constructed.

*Date format: YYYY/MM/DD